<u>Theorem</u>: For task *i*, the minimum nonzero separation of a task instance release time from the corresponding frame start time is  $gcd(f, T_i)$ .

## Proof:

Let  $g = \text{gcd}(f, T_i)$ , it follows from the definition of the greatest common divisor (gcd) that g must divide each of f and  $T_i$ . Let i be a task with zero phasing. Now, assume that this theorem is violated for certain integers m and n, such that the nth release of task i occurs in the mth frame and the difference of the start of the mth frame and the release of the nth instance is a nonzero value less than g, that is

$$0 < (m \times f - n \times T_i) < g$$

Dividing this expression throughout by g, we get

$$0 < (m \times f/g - n \times T_i/g) < 1$$

However, both f/g and  $T_i/g$  are integers because  $g = \text{gcd}(f, T_i)$ . Therefore we can write  $f/g = I_1$ , and  $T_i/g = I_2$  for some integer values  $I_1$  and  $I_2$ . Substituting in the above inequality we get

$$0 < (m \times I_1 - n \times I_2) < 1$$

Since  $m \times I_1$  and  $n \times I_2$  are both integers, their difference cannot be a fractional value between 0 and 1. Therefore this inequality can never be satisfied.

It can therefore be concluded that the minimum time between a frame boundary and the arrival of the corresponding instance of task cannot be less than  $gcd(f, T_i)$ . [Actually it cannot be less than any common divisor, thus its minimum value is the gcd.]

(Proof adapted from notes by Rajib Mall).