(a) Account A can retrieve or modify any relation except DEPENDENT and can grant any of these privileges to other users.

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	-----	-------	---------	-----	--------	-----------	-----

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname	Pnumber	Plocation	Dnum
-------	---------	-----------	------

#### WORKS\_ON

Essn	Pno	Hours
------	-----	-------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

(a) Account A can retrieve or modify any relation except DEPENDENT and can grant any of these privileges to other users.

### GRANT SELECT, UPDATE

ON EMPLOYEE, DEPARTMENT, DEPT\_LOCATIONS, PROJECT, WORKS\_ON TO USER\_A WITH GRANT OPTION;

Enomo	Minit	Lanna	Sen	Delata	Address	Cou	Colore	Super_ssn	Dee
rname	IVIIIII	Lname	OSI	Duate	Address	Sex	Salary	Super_ssn	Dho

#### DEPARTMENT

EMPLOYEE

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname	Pnumber	Plocation	Dnum
-------	---------	-----------	------

#### WORKS\_ON

Feen	Pno	Hours
Lasin	FIIU	Hours

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

(b) Account B can retrieve all the attributes of EMPLOYEE and DEPARTMENT

except for SALARY, MGRSSN, and MGRSTARTDATE.

MPLOYE	E								
Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno

#### DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname	Pnumber	Plocation	Dnum
-------	---------	-----------	------

#### WORKS\_ON

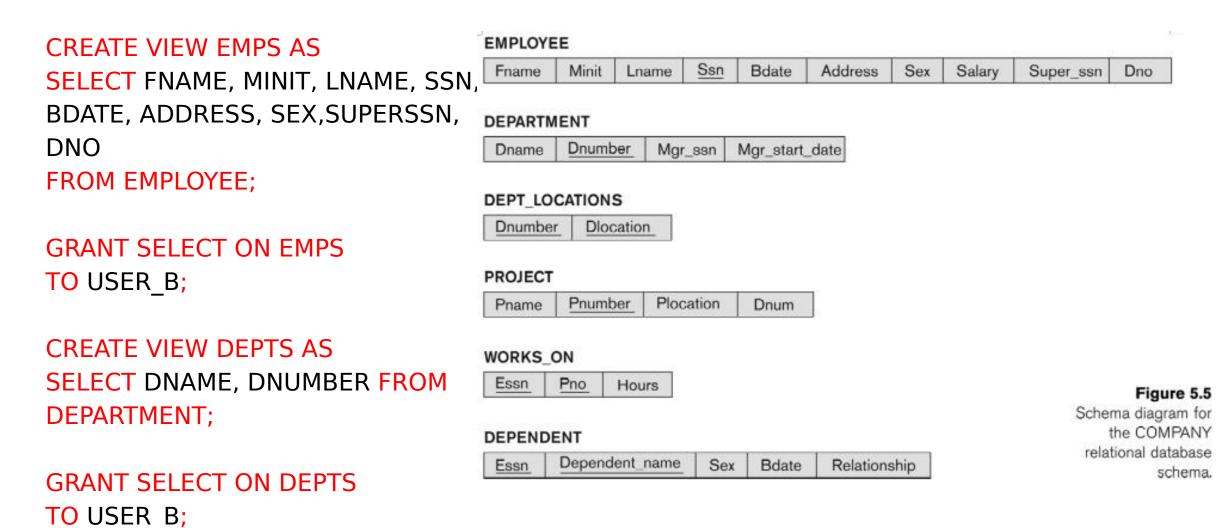
Essn	Pno	Hours
------	-----	-------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

23.24 Consider the relational database schema Write SQL statements to grant these privileges. Use views were appropriate:(b) Account B can retrieve all the attributes of EMPLOYEE and DEPARTMENT except for SALARY, MGRSSN, and MGRSTARTDATE.



(c) Account C can retrieve or modify WORKS\_ON but can only retrieve the

FNAME, MINIT, LNAME, SSN attributes of EMPLOYEE and the PNAME,

PNUMBER attributes of PROJECT.

1	Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno	
---	-------	-------	-------	-----	-------	---------	-----	--------	-----------	-----	--

#### DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname Pnun	nber Plocat	ion Dnum
------------	-------------	----------

#### WORKS\_ON

Essn	Pno	Hours
------	-----	-------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

(c) Account C can retrieve or modify WORKS\_ON but can only retrieve the

FNAME, MINIT, LNAME, SSN attributes of EMPLOYEE and the PNAME,

PNUMBER attributes of PROJECT.

GRANT SELECT, UPDATE ON WORKS	ON	DEPA
		Door

TO USER\_C;

CREATE VIEW EMP1 AS SELECT FNAME, MINIT, LNAME, SSN FROM EMPLOYEE;

GRANT SELECT ON EMP1 TO USER\_C;

	EMPLOYE	EE					
	Fname	Minit	Lname	Ssn	Bdate	Address	Sex
RKS ON		AENT					
	DEPARTN	ENI					

#### DEPT\_LOCATIONS



#### PROJECT

	Pname	Pnumber	Plocation	Dnum
--	-------	---------	-----------	------

#### WORKS\_ON

Essn Pno Hours
----------------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure

Schema diagram the COMPA relational databa scher

Super\_ssn

Dno

Salary

(c) Account C can retrieve or modify WORKS\_ON but can only retrieve the

FNAME, MINIT, LNAME, SSN attributes of EMPLOYEE and the PNAME,

PNUMBER attributes of PROJECT.

CREATE VIEW PROJ1 AS SELECT PNAME, PNUMBER FROM PROJECT;

GRANT SELECT ON PROJ1 TO USER\_C:

MPLOYE	E								
Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno

#### DEPARTMENT

Dhame Dhumber Mgr_ssh Mgr_start_date	Dname	Dnumber	Mgr_ssn	Mgr_start_date
--------------------------------------	-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber	Dlocation
---------	-----------

#### PROJECT

Thanke Dian	Pname	Pnumber	Plocation	Dnum
-------------	-------	---------	-----------	------

#### WORKS\_ON

|--|

#### DEPENDENT



#### Figure !

Schema diagram the COMPA relational databa scher

(d) Account D can retrieve any attribute of EMPLOYEE or DEPENDENT and can modify DEPENDENT.

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	-----	-------	---------	-----	--------	-----------	-----

#### DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname	Pnumber	Plocation	Dnum
-------	---------	-----------	------

#### WORKS\_ON

Essn	Pno	Hours
------	-----	-------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

(d) Account D can retrieve any attribute of EMPLOYEE or DEPENDENT and can modify DEPENDENT.

### GRANT SELECT ON EMPLOYEE, DEPENDENT TO USER\_D;

GRANT UPDATE ON DEPENDENT TO USER\_D;

	1.		1000 C 201			11.122	NOTABLE INC.		
Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno

#### DEPARTMENT

EMPLOYEE

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname	Pnumber	Plocation	Dnum
-------	---------	-----------	------

#### WORKS\_ON

Essn	Pno	Hours
------	-----	-------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

(e) Account E can retrieve any attribute of EMPLOYEE but only for EMPLOYEE tuples that have DNO = 3.

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
-------	-------	-------	-----	-------	---------	-----	--------	-----------	-----

#### DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
-------	---------	---------	----------------

#### DEPT\_LOCATIONS

Dnumber Dlocation

#### PROJECT

Pname	Pnumber	Plocation	Dnum
-------	---------	-----------	------

#### WORKS\_ON

Essn	Pno	Hours
------	-----	-------

#### DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
------	----------------	-----	-------	--------------

#### Figure 5.5

(e) Account E can retrieve any attribute of EMPLOYEE but only for EMPLOYEE tuples that have DNO = 3.

----

	EMPLOY	EE								
	Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
CREATE VIEW DNO3 EMPLOYEES	5									
—	DEPART	IENT								
AS SELECT * FROM EMPLOYEE	Dname	Dnumb	er Mg	_ssn	Mgr_start	_date				
WHERE DNO=3;										
	DEPT_LO									
	Dnumbe	r Dloc	ation							
GRANT SELECT ON	PROJECT									
DNO2 EMDLOVEES TO LISED E.	Pname	Pnumb	er Plor	ation	Dnum	1				
DNO3_EMPLOYEES TO USER_E;	1 Hame			Auton	Diam					
	WORKS	ON								
	Essn	_	Hours							El auros d
									Sche	Figure ! ma diagram
	DEPEND	ENT							1	the COMPAI
	Essn	Depend	ent_name	Sex	Bdate	Relation	ship		rela	tional databa scher

# How would it appear to a user with classification U?

### (d) EMPLOYEE

Name	Salary	JobPerformance	TC
Smith U	40000 C	Fair S	S
Smith U	40000 C	Excellent C	C
Brown C	80000 S	Good C	S

# How would it appear to a user with classification U?

### (d) EMPLOYEE

Name	Salary	JobPerformance	TC
Smith U	40000 C	Fair S	S
Smith U	40000 C	Excellent C	C
Brown C	80000 S	Good C	S

Name	Salary	JobPerformance	ТС
Smith U	NULL	NULL	NULL

Suppose a classification U user tries to update the salary of "Smith" to \$50,000; what would be the result of this action?

#### (d) EMPLOYEE

Name	Salary	JobPerformance	TC
Smith U	40000 C	Fair S	S
Smith U	40000 C	Excellent C	C
Brown C	80000 S	Good C	S

Name	Salary	JobPerformance	ТС
Smith U	40000 U	Fair S	S
Smith U	40000 U	Excellent C	С
Smith U	50000 U	NULL	U
Brown C	80000 S	Good C	S

- Incremental logging with deferred updates implies that the recovery system must necessarily
- a. store the old value of the updated item in the log.
- b. store the new value of the updated item in the log.
- c. store both the old and new value of the updated item in the log.
- d. store only the Begin Transaction and Commit Transaction records in the log.

- Incremental logging with deferred updates implies that the recovery system must necessarily
- a. store the old value of the updated item in the log.

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c. store both the old and new value of the updated item in the log.

d. store only the Begin Transaction and Commit Transaction records in the log.

The write-ahead logging (WAL) protocol simply means that

a. writing of a data item should be done ahead of any logging operation.

b. the log record for an operation should be written before the actual data is written.

- c. all log records should be written before a new transaction begins execution.
- d. the log never needs to be written to disk.

The write-ahead logging (WAL) protocol simply means that

a. writing of a data item should be done ahead of any logging operation.

## b. the log record for an operation should be written before the actual data is written.

c. all log records should be written before a new transaction begins execution.

d. the log never needs to be written to disk.

In case of transaction failure under a deferred update incremental logging scheme, which of the following will be needed?

- a. an undo operation
- b. a redo operation
- c. an undo and redo operation
- d. none of the above

In case of transaction failure under a deferred update incremental logging scheme, which of the following will be needed?

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### **b.** a redo operation

- c. an undo and redo operation
- d. none of the above

For incremental logging with immediate updates, a log record for a transaction would contain

a. a transaction name, a data item name, and the old and new value of the item.

b. a transaction name, a data item name, and the old value of the item.

- c. a transaction name, a data item name, and the new value of the item.
- d. a transaction name and a data item name.

For incremental logging with immediate updates, a log record for a transaction would contain

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- d. a transaction name and a data item name.

## For correct behavior during recovery, undo and redo operations must be

- a. commutative.
- b. associative.
- c. idempotent.
- d. distributive.

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## c. idempotent.

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There is a possibility of a cascading rollback when

a. a transaction writes items that have been written only by a committed transaction.

b. a transaction writes an item that is previously written by an uncommitted transaction.

c. a transaction reads an item that is previously written by an uncommitted transaction.

d. both (b) and (c).

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d. both (b) and (c).

- To cope with media (disk) failures, it is necessary
  - a. for the DBMS to only execute transactions in a single user environment.
  - b. to keep a redundant copy of the database.
  - c. to never abort a transaction.
  - d. all of the above.

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a. for the DBMS to only execute transactions in a single user environment.

### **b.** to keep a redundant copy of the database.

c. to never abort a transaction.

d. all of the above.

## If the shadowing approach is used for flushing a data item back to disk, then

- a. the item is written to disk only after the transaction commits.
- b. the item is written to a different location on disk.
- c. the item is written to disk before the transaction commits. d. the item is written to the same disk location from which it was read.

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a. the item is written to disk only after the transaction commits.

## **b.** the item is written to a different location on disk.

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Aries

For recovery in ARIES, study the shown log at time of crash.

- (a) Show the transaction table and the dirty page table at time of check point and after the analysis phase of the recovery process.
- (b) Explain what will happen in the other two phases of the recovery process

LSN	$Last\_LSN$	Trans_Id	Type	$Page_Id$
1	0	T1	Update	Α
2	0	Τ2	Update	В
3	2	T2	Commit	
4	Begin checkpoint			
5	End checkpoint			
6	1	T1	Commit	$\mathbf{C}$
7	0	T3	Update	
8	7	Τ3	Commit	Α